NEUROSCIENCE IN THE DATABASE

MAPPING THE MORPHOLOGY OF NEURONS AND SYNAPSES WITH CATMAID AND POSTGIS

Tom Kazimiers | kazmos GmbH

WHO?



- Tom Kazimiers, Main developer of CATMAID, MSc CS (Dipl.Inf.) TU Dresden
- 3+ years at MPI CBG (Dresden, DE)
- 6+ years at HHMI Janelia Research Campus (Ashburn, Virginia, USA)
- Since 2020: Open Source Research Software Engineer and founder of consultancy kazmos GmbH (DE)

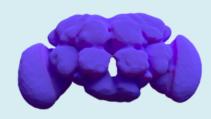
TOPICS

- Motivation and data
- Overview
- Data representation
- Neuron graph queries
- PostGIS queries

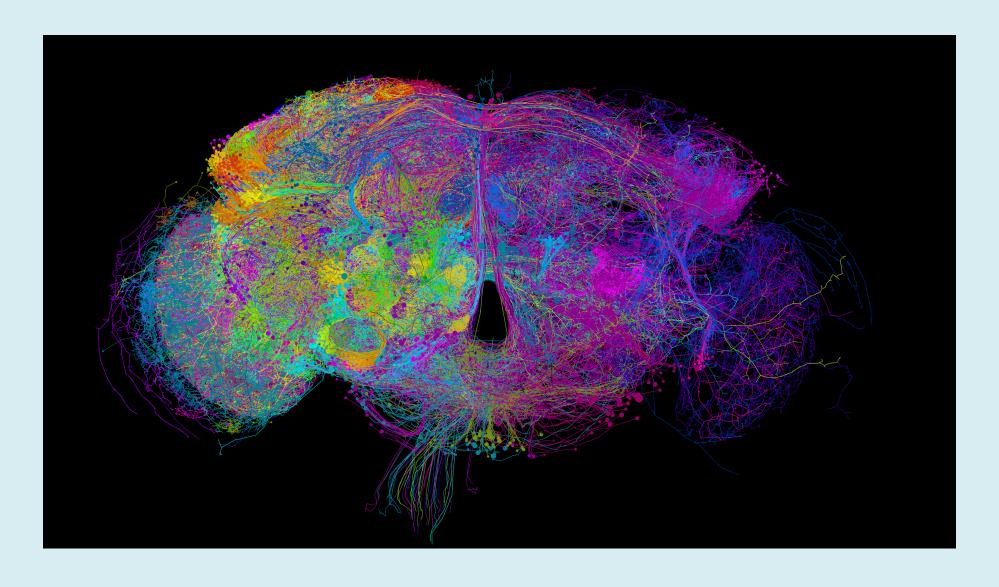
MOTIVATION: CONNECTOMICS

- Structure/development of central nervous system
- Function of specific neurons and neuron classes
- Synapses and resulting networks
- How is behavior controlled?
- Fundamental research



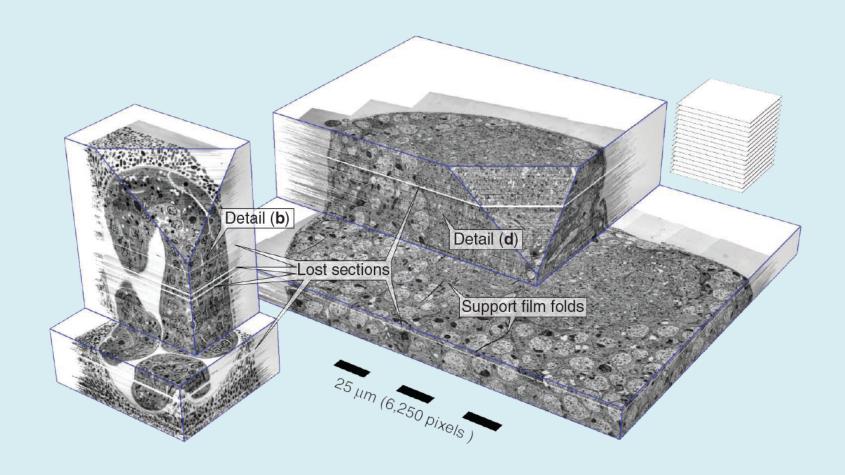


Model organism *Drosophila Melanogaster* ("fruit fly") Images: Wikimedia, CC-BY-SA 2.5 (left) and CC-BY-SA 4.0 (right)



~7 meters of neurons in brain of fruit fly, 20 labs, 186 users created and reviewed 33 million nodes in more than 60 person years Image by Philipp Schlegel, Jefferis lab

IMAGE DATA: ELECTRON MICROSCOPY



ssTEM: Elastic volume reconstruction from series of ultra-thin microscopy sections S. Saalfeld, R. Fetter, A. Cardona, et al., Nature Methods, 2012

REPRESENTATIVE DATA SETS

- Drosophila Larva L1 Dataset (ssTEM):
 28,128 x 31,840 x 4,841 px @ 3.8 x 3.8 x 50 nm/px
 0.9TB as JPEG tiles (512 x 512)
- Drosophila FAFB Dataset (ssTEM):
 253,952 x 155,648 x 7,063 px @ 4 x 4 x 40 nm/px
 10.9TB as JPEG tiles (1024 x 1024)
- Drosophila Hemibrain Dataset (FIBSEM):
 40,959 x 34,815 x 43,007 px @ 8 x 8 x 8 nm/px
 5.3TB as Neuroglancer Precomputed (64 x 64 x 64)
- Compare: OpenStreetMap tiles (z18): Scandinavian countries: 1TB (Geofabrik)

North + South America: 11TB (Geofabrik)

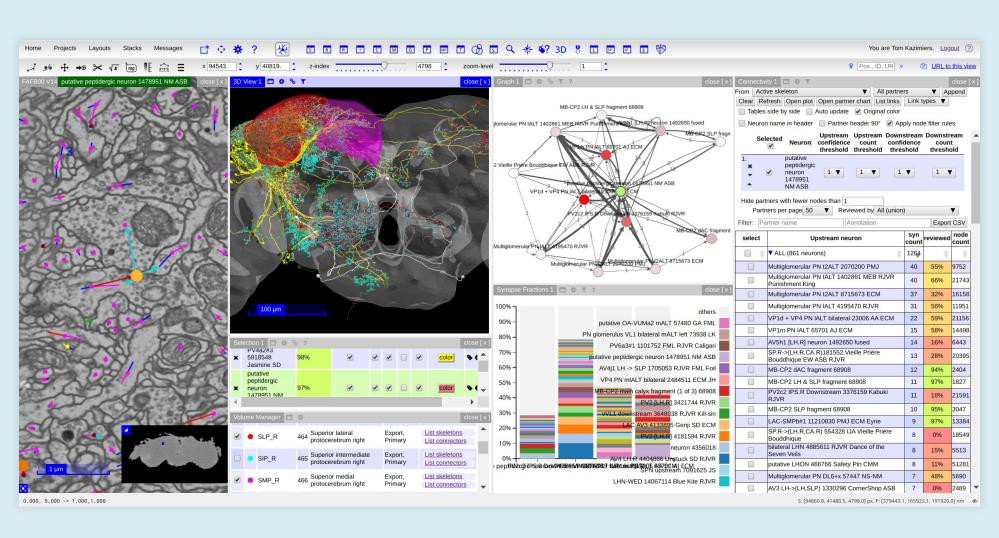
North America: 5.4TB (Geofabrik)

CATMAID

Collaborative Annotation Toolkit for

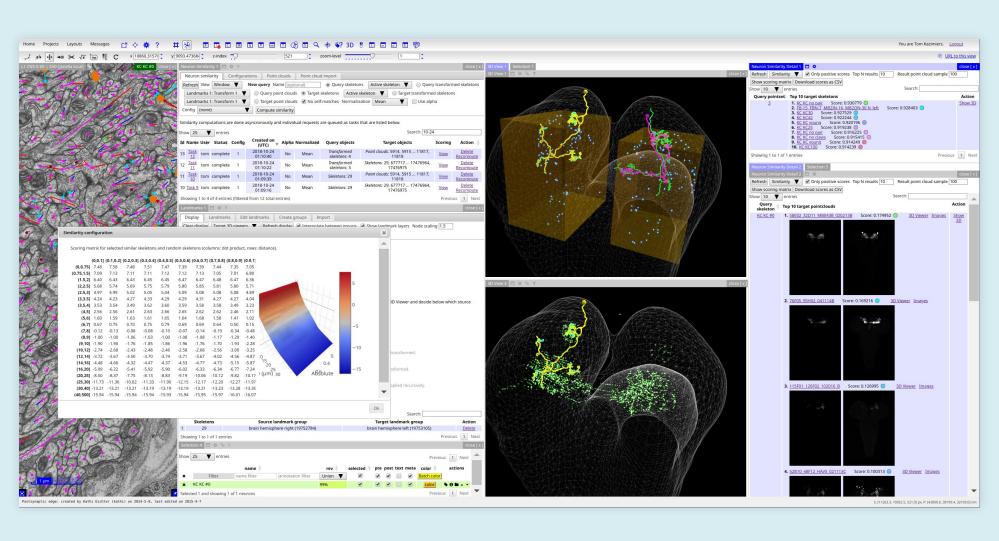
Massive Amounts of Image Data

USER INTERFACE I



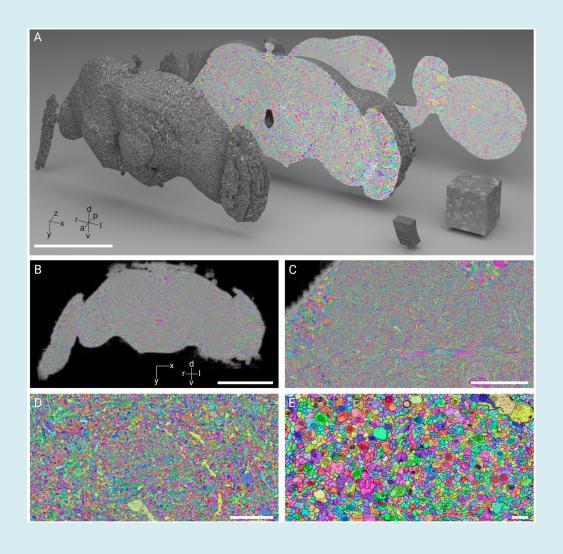
Typical CATMAID workspace with 2D and 3D views plus some connectivity tools Example circuit from Dolan et al. 2019, FAFBv14 dataset

USER INTERFACE II



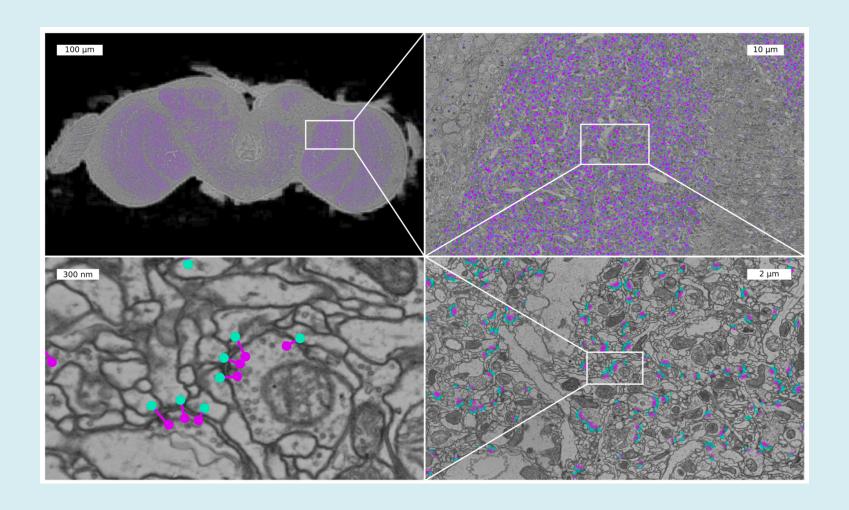
CATMAID with neuron similarity tools and point-clouds from light data Screenshots from Albert Cardona's Drosophila L1 dataset

NEURON SEGMENTATION



Segmentation of FAFB dataset: used also for skeletonization; Similar in FlyWire project Li, Jain et al, bioRxiv 2020

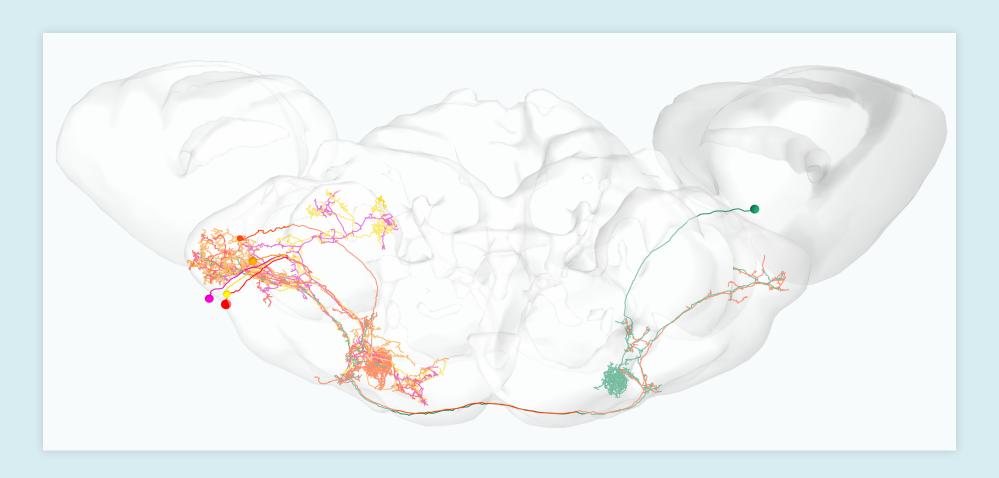
SYNAPTIC PARTNER DETECTION



Automatic detection of synaptic partners (CircuitMap tool in CATMAID)

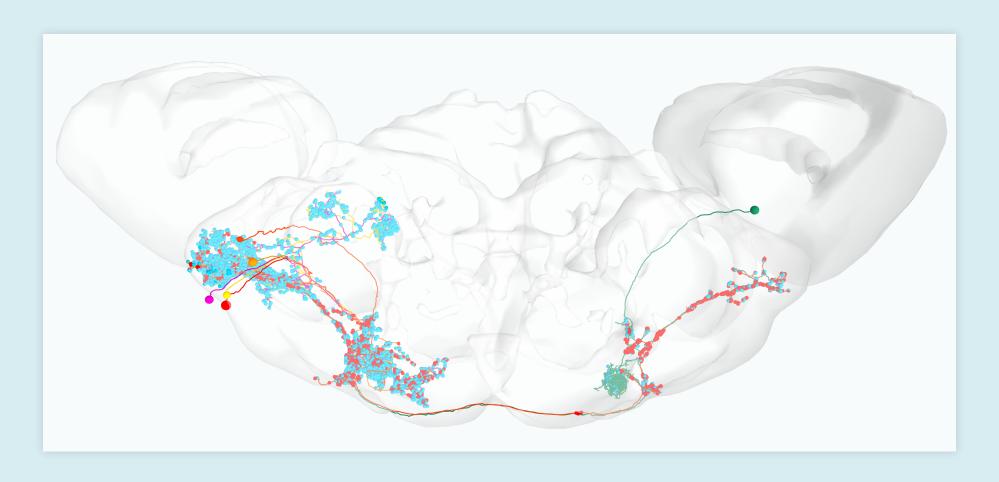
Buhman et al, Nature Methods, 2021

MODELLING NEURONS



First order (downstream) partners (orange/red) to MBON a2sc Left ASB (green) From Dolan and Belliart-Guérin et al. 2018

MODELLING NEURONS

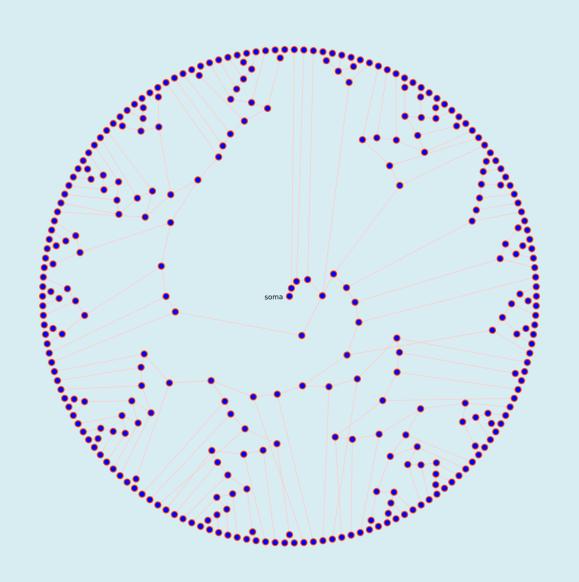


MBON and PD neurons with synapses highlighted (pre: red, post: cyan) From Dolan and Belliart-Guérin et al. 2018

MORPHOLOGY OF A NEURON



TOPOLOGY OF A NEURON



REPRESENTING NEURONS IN A DATABASE

- Relational data, explicit schema: PostgreSQL
- Nodes of all skeletons in single table
- Every element knows parent or is root (trees)

```
\d treenode
    Column
                                    Nullable
                                                Default
                      Type
id
                 bigint
                                    not null
parent_id
                 bigint
skeleton_id
                 bigint
                                    not null
project_id
                                    not null
                 integer
 location_x
                 real
                                    not null
Foreign-key constraints:
 "treenode_parent_id_fkey" FOREIGN KEY (parent_id)
                            REFERENCES treenode(id)
```

QUERY NODES OF A SKELETON

Query ignores any structure/relation between nodes

```
1 SELECT id, parent_id, location_x, location_y, location_z
2 FROM treenode
3 WHERE skeleton_id = {skeleton_id}
```

Root node of skeleton

```
1 SELECT id, parent_id, location_x, location_y, location_z
2 FROM treenode
3 WHERE skeleton_id = {skeleton_id} AND parent_id IS NULL
```

QUERY NEIGHBOR NODES

```
1 SELECT c.id, c.location_x, c.location_y, c.location_z
2 FROM treenode c
3 WHERE parent_id = {node_id}
4
5 UNION ALL
6
7 SELECT p.id, p.location_x, p.location_y, p.location_z
8 FROM treenode t
9 JOIN treenode p
10 ON p.id = t.parent_id
11 WHERE t.id = {node_id}
```

```
1 WITH RECURSIVE nodes (id, depth) AS (
2    SELECT t.id, 1
3    FROM treenode t
4    WHERE t.parent_id IS NULL
5     AND t.skeleton_id = {skeleton_id}
6    UNION ALL
7    SELECT t.id, p.depth + 1
8    FROM treenode t
9    JOIN nodes p ON t.parent_id = p.id
10 )
11 SELECT t.id, t.parent_id
12 FROM treenode t
13 WHERE t.skeleton_id = {skeleton_id}
14 AND NOT EXISTS (SELECT n.id FROM nodes n WHERE n.id = t.id);
```

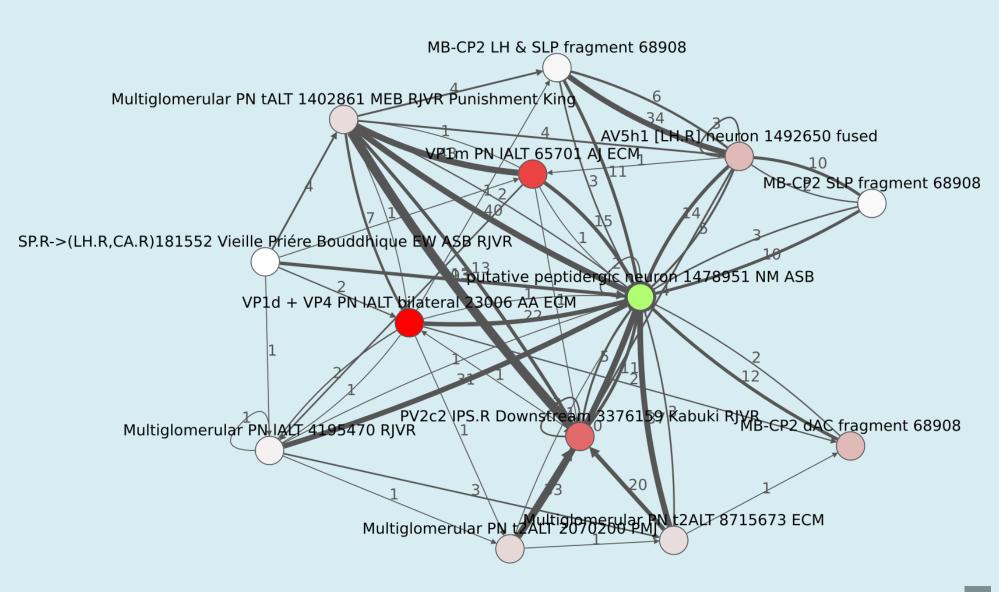
```
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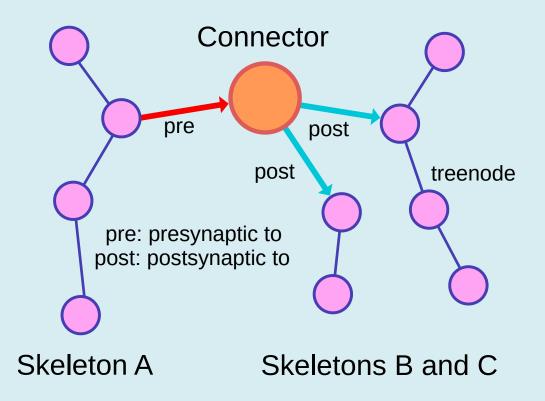
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```

CONNECTED NEURONS



NEURON GRAPHS IN CATMAID



NEURON GRAPHS IN SQL

```
# \d treenode_connector
      Column
                                      Nullable |
                                                 Default
                         Type
   relation_id
                   bigint
                                      not null
   treenode_id
                   bigint
                                      not null
   connector_id
                   bigint
                                      not null
                   bigint
   skeleton_id
                                      not null
   project_id
                   integer
                                      not null
   confidence
                   smallint
                                      not null |
                                                5
  Foreign-key constraints:
```

- Useful for direct connections between skeletons, e.g. find partner skeleton IDs and synapse counts
- Gets more complex with additional levels

```
1 SELECT tc2.skeleton_id AS partner, COUNT(*) AS n_post_sites
2 FROM treenode_conector tc1
3 JOIN treenode_conector tc2
4    ON tc1.connector_id = tc2.connector_id
5    AND tc1.treenode_id != tc2.treenode_id
6 WHERE tc1.relation_id = {presynaptic_to_id}
7    AND tc2.relation_id = {postsynaptic_to_id}
8    AND tc1.project_id = {project_id}
9    AND tc2.project_id = {project_id}
10    AND tc1.skeleton_id = {skeleton_id}
11 GROUP BY tc2.skeleton_id
```

- Useful for direct connections between skeletons, e.g. find partner skeleton IDs and synapse counts
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```
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5    AND tc1.treenode_id != tc2.treenode_id
6 WHERE tc1.relation_id = {presynaptic_to_id}
7    AND tc2.relation_id = {postsynaptic_to_id}
8    AND tc1.project_id = {project_id}
9    AND tc2.project_id = {project_id}
10    AND tc1.skeleton_id = {skeleton_id}
11 GROUP BY tc2.skeleton_id
```

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7    AND tc2.relation_id = {postsynaptic_to_id}
8    AND tc1.project_id = {project_id}
9    AND tc2.project_id = {project_id}
10    AND tc1.skeleton_id = {skeleton_id}
11 GROUP BY tc2.skeleton_id
```

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6 WHERE tc1.relation_id = {presynaptic_to_id}
7   AND tc2.relation_id = {postsynaptic_to_id}
8   AND tc1.project_id = {project_id}
9   AND tc2.project_id = {project_id}
10   AND tc1.skeleton_id = {skeleton_id}
11 GROUP BY tc2.skeleton_id
```

- Useful for direct connections between skeletons, e.g. find partner skeleton IDs and synapse counts
- Gets more complex with additional levels

```
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5   AND tc1.treenode_id != tc2.treenode_id
6 WHERE tc1.relation_id = {presynaptic_to_id}
7   AND tc2.relation_id = {postsynaptic_to_id}
8   AND tc1.project_id = {project_id}
9   AND tc2.project_id = {project_id}
10   AND tc1.skeleton_id = {skeleton_id}
11 GROUP BY tc2.skeleton_id
```

POSTGIS DATA

- In 2D view: intersections of skeletons with 2D plane and store edges between nodes
- Main planar access in a single direction (usually Z)

```
# \d treenode_edge
     Column
                        Type
                                      | Collation | Nullable | Default
   id | bigint
                                                    not null
   parent_id | bigint
   project_id | integer
                                                    not null
            | geometry(LineStringZ)
   edge
                                                    not null |
  Indexes:
      "treenode_edge_pkey" PRIMARY KEY, btree (id)
      "treenode_edge_project_id_index" btree (project_id)
      "treenode_edge_2d_gist" gist (edge)
      "treenode_edge_3d_gist" gist (edge gist_geometry_ops_nd)
      "treenode_edge_z_range_gist" gist (floatrange(
          st_zmin(edge::box3d), st_zmax(edge::box3d), '[]'::text))
```

BOUNDING BOX QUERIES

- 2D view: cross section of all neurons in a field of view at a particular depth
- Different FoV BBs can benefit from different indexes, so CATMAID tries to use fitting query, e.g.:
 - Whole slice is visible
 - Medium-sized rectangle
 - Small FoV with BB closer to a cube
- For analysis/filters: intersect with 3D meshes (e.g. compartments)

LARGE FIELD OF VIEW

```
1 SELECT *
 2 FROM (
     SELECT DISTINCT ON (id) UNNEST(ARRAY[te.id, te.parent_id]) AS id
     FROM treenode edge te
     WHERE te.project id = {project id}
     AND floatrange(ST_ZMin(te.edge), ST_ZMax(te.edge), '[]')
         && floatrange({z1}, {z2}, '[)')
     AND te.edge && ST_MakeEnvelope({left}, {top}, {right}, {bottom})
     AND ST_3DDWithin(te.edge, ST_MakePolygon(ST_MakeLine(ARRAY)
         ST_MakePoint({left}, {top}, {halfz}),
         ST_MakePoint({right}, {top}, {halfz}),
11
12
         ST_MakePoint({right}, {bottom}, {halfz}),
13
         ST_MakePoint({left}, {bottom}, {halfz}),
14
         ST_MakePoint({left}, {top}, {halfz})]::geometry[])),
         {halfzdiff})
15
    bb treenode
16
17 JOIN treenode t1
     ON t1.id = bb_treenode.id
18
```

Get all child and parent nodes of edges in a 3D bounding box

LARGE FIELD OF VIEW

```
AND floatrange(ST_ZMin(te.edge), ST_ZMax(te.edge), '[]')
    && floatrange({z1}, {z2}, '[)')
```

```
Constrain result edges to those in a user-defined Z range and allow use of index "treenode_edge_z_range_gist" gist (floatrange(ST_ZMin(edge::box3d), ST_ZMax(edge::box3d), '[]'::text))
```

LARGE FIELD OF VIEW

```
AND te.edge && ST_MakeEnvelope({left}, {top}, {right}, {bottom})
```

Constrain result edges to only those in XY user-defined XY area and allow use of index "treenode_edge_2d_gist" gist (edge);

LARGE FIELD OF VIEW

```
AND ST_3DDWithin(te.edge, ST_MakePolygon(ST_MakeLine(ARRAY)
         ST_MakePoint({left}, {top}, {halfz}),
         ST_MakePoint({right}, {top}, {halfz}),
11
12
         ST_MakePoint({right}, {bottom}, {halfz}),
13
         ST_MakePoint({left}, {bottom}, {halfz}),
14
         ST_MakePoint({left}, {top}, {halfz})]::geometry[])),
         {halfzdiff})
15
```

Test true distance of edge to BB to remove matches where only the BB of the edge is close, Lowers *false positives* where only query BB and edge BB intersect

SMALL FOV / CUBE-LIKE BB

```
1 SELECT *
 2 FROM (
     SELECT DISTINCT ON (id) UNNEST(ARRAY[te.id, te.parent_id]) AS id
    FROM treenode edge te
     WHERE te.project id = {project id}
     AND te.edge &&& ST MakeLine(ARRAY[
         ST_MakePoint({left}, {bottom}, {z2}),
         ST_MakePoint({right}, {top}, {z1})] ::geometry[])
     AND ST_3DDWithin(te.edge, ST_MakePolygon(ST_MakeLine(ARRAY[
         ST_MakePoint({left}, {top}, {halfz}),
         ST_MakePoint({right}, {top}, {halfz}),
11
         ST_MakePoint({right}, {bottom}, {halfz}),
12
         ST_MakePoint({left}, {bottom}, {halfz}),
13
14
         ST_MakePoint({left}, {top}, {halfz})]::geometry[])),
         {halfzdiff})
15
    bb treenode
17 JOIN treenode t1
     ON t1.id = bb_treenode.
18
```

Similar structure like other query, but different index Small result set, no extra filter conditions are beneficial

SMALL FOV / CUBE-LIKE BB

```
AND te.edge &&& ST MakeLine(ARRAY[
    ST_MakePoint({left}, {bottom}, {z2}),
    ST_MakePoint({right}, {top}, {z1})] ::geometry[])
```

The &&& operator allows the planner to use the index "treenode_edge_3d_gist" gist (edge gist_geometry_ops_nd)

SMALL FOV / CUBE-LIKE BB

```
AND ST_3DDWithin(te.edge, ST_MakePolygon(ST_MakeLine(ARRAY)
         ST_MakePoint({left}, {top}, {halfz}),
         ST_MakePoint({right}, {top}, {halfz}),
11
12
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         ST_MakePoint({left}, {bottom}, {halfz}),
14
         ST_MakePoint({left}, {top}, {halfz})]::geometry[])),
         {halfzdiff})
15
```

Limit the number of false positive BB-only intersections, just like before

CLOSEST NODE IN 3D

```
1 SELECT treenode.id, skeleton_id, location_x, location_y, location_z
2 FROM treenode
3 JOIN (
4     SELECT id, edge
5     FROM treenode_edge
6     WHERE project_id = {project_id}
7     ORDER BY edge <<->> ST_MakePoint({x}, {y}, {z})
8     LIMIT 100
9 ) closest_node(id, edge)
10 ON closest_node.id = treenode.id
11 ORDER BY ST_StartPoint(edge) <<->> ST_MakePoint({x}, {y}, {z})
12 LIMIT 1
```

Find closest node among the 100 closest edge centeroids in order to allow 3D index use. Assume closest node is among closest edge BB centroids (<<->> operator)

CLOSEST NODE IN 3D

```
1 SELECT treenode.id, skeleton_id, location_x, location_y, location_z
2 FROM treenode
3 JOIN (
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10 ON closest_node.id = treenode.id
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12 LIMIT 1
```

Find 100 closest (euclidean) edges

CLOSEST NODE IN 3D

```
1 SELECT treenode.id, skeleton_id, location_x, location_y, location_z
2 FROM treenode
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11 ORDER BY ST_StartPoint(edge) <<->> ST_MakePoint({x}, {y}, {z})
12 LIMIT 1
```

Get closest edge start node of closest edges

MORE POSTGIS

- More spatial annotations (POIs, bookmarks, text)
- Storing 3D meshes as TIN data
- Store spatial meta data for point clouds (BB, center of mass)

Thank you PostGIS team for making spatial data in the DB easier, even beyond GIS!

ACKNOWLEDGMENTS

CONTRIBUTORS

- Albert Cardona
- Andrew Champion
- Chris Barnes
- Stephan Gerhard
- Davi Bock
- Stephan Saalfeld
- Marta Costa
- Greg Jefferis
- Will Patton
- Philipp Schlegel
- Alex Bates

DATA CREATION AND DATA WRANGLING

- Cardona Lab Tracers
- Wellcome Trust Tracers
- CAT Team @ JRC
- Eric Perlman @ Yikes, LLC

LABS AND TEAMS

- Bock Lab @ UVM
- Cardona Lab @ LMB
- Lippincott-Schwartz Lab @ JRC
- Jekely Lab @ UEX
- Lee Lab @ HMS
- Funke Lab @ JRC

QUESTIONS?

- catmaid.org
- github.com/catmaid/catmaid
- tom@kazmos.de
- @tomkazimiers
- Docker image:

docker run -p 8080:80 catmaid/catmaid